



A Simulation-Based Approach to Differential Privacy

Younes Boulaguiem, PhD
Postdoctoral Researcher in Statistics

University of Geneva – Faculty of Medicine
Clinical Research Center, HUG

A simulation-based fiducial matching approach for DP inference.
Delivered May 2023

[View Preprint](#)

Motivating Problem

- Suppose we are implementing a survey to gather some **sensitive information** from a population of n individuals.
- For example we are interested to know if the proportion of individuals that have committed a traffic light violation within a time frame. We denote the information of each individual i , $X_i \in \{0, 1\}$.
- The analyst then gathers each individual response Y_i which may or may not correspond to the true value X_i :

$$Y_i = \begin{cases} X_i & \text{with probability } 1/2 + \delta \\ 1 - X_i & \text{with probability } 1/2 - \delta \end{cases}$$

- The analyst is interested in getting an estimate $\tilde{p} = \frac{1}{n} \sum_{i=1}^n Y_i$ of $p = \frac{1}{n} \sum_{i=1}^n X_i$.

The Non-private and the Perfectly Private frameworks

$$Y_i = \begin{cases} X_i & \text{with probability } 1/2 + \delta \\ 1 - X_i & \text{with probability } 1/2 - \delta \end{cases}$$

- In the **non-private** framework $\delta = 1/2$ and p can be **retrieved precisely** with \tilde{p} as $\tilde{p} = \frac{1}{n} \sum_{i=1}^n X_i = p$. However, the analyst has **complete access** to the private information X . There is no privacy.
- In the **perfectly private** framework, $\delta = 0$ and the distribution of Y_i is a uniform and is **completely independent** of X_i . We have complete privacy, but on the flip side, \tilde{p} is **completely uninformative** about p .

The Randomized Response Algorithm

$$Y_i = \begin{cases} X_i & \text{with probability } 1/2 + \delta \\ 1 - X_i & \text{with probability } 1/2 - \delta \end{cases}$$

- Let's consider an intermediate strategy, and put $\delta = 1/4$. By observing $Y_i = X_i$ with probability $3/4$, meaning that the individual's true information could still be $1 - Y_i$ with probability $1/4$, it is said that the individual has been granted **“plausible deniability”**.
- How “deniable” the response is corresponds to **the level of privacy** they have been afforded, i.e., **the smaller the δ** , the greater the deniability, or in other words, **the greater the privacy**.
- Moreover, defining $\tilde{p} = \frac{1}{n} \sum_{i=1}^n \frac{1}{2\delta} (Y_i + \delta - 1/2)$, it follows that $|\tilde{p} - p| = O_p(\frac{1}{\delta\sqrt{n}})$.

Central Differential Privacy

- The previous example falls under the more general framework of Differential Privacy (DP), a broad formalization of the notion of “plausible deniability”.
- We consider in the following a version of DP called central DP.
- It involves a trusted curator who receives raw data from individuals, runs a randomized algorithm and publicly outputs its results.
- DP is a property of this algorithm stating that no individual data has a significant impact on the output of the algorithm.

ϵ -Differential Privacy

ϵ -Differential Privacy

Consider any two neighbouring datasets $x, y \in \mathcal{X}^n$ that differ in exactly one entry, which we denote $x \sim y$.

A randomized algorithm $\mathcal{M} : \mathcal{X}^n \rightarrow \mathcal{Y}$ is ϵ -differentially private if for all neighbouring datasets x, y , and any event $S \subseteq \mathcal{Y}$

$$\Pr(\mathcal{M}(x) \in S) \leq e^\epsilon \Pr(\mathcal{M}(y) \in S)$$

where ϵ is called the **privacy budget** and the randomness is over the choices of \mathcal{M} .

When ϵ is small, $e^\epsilon \approx 1 + \epsilon$ and the inequality states that the amount of information one would learn about a given individual would not change (much) whether or not they are in the dataset.

Post-processing Property

Let $\mathcal{M} : \mathcal{X}^n \rightarrow \mathcal{Y}$ be ϵ -differentially private, and let $f : \mathcal{Y} \rightarrow \mathcal{Y}'$ be an arbitrary randomized mapping. Then $f \circ \mathcal{M}$ is ϵ -differentially private.

The Laplace Mechanism I

ℓ_1 -sensitivity

Let $f : \mathcal{X}^n \rightarrow \mathbb{R}^k$. The ℓ_1 -sensitivity of f is

$$\Delta f = \max_{X \sim X'} \|f(X) - f(X')\|_1$$

The Laplace Distribution

The Laplace distribution with location and scale parameters 0 and b , respectively, has the following density:

$$p(x) = \frac{1}{2b} \exp\left(-\frac{|x|}{b}\right).$$

We note $Y \sim \text{Lap}(b)$, with $\mathbb{E}[Y] = 0$ and $\text{var}(Y) = 2b^2$. Furthermore, the following holds:

$$\Pr(|Y| \geq tb) = \exp(-t).$$

The Laplace Mechanism II

The Laplace Mechanism

Let $f : \mathcal{X}^n \rightarrow \mathbb{R}^k$. The Laplace mechanism is defined as

$$\mathcal{M}(X) = f(X) + (Y_1, \dots, Y_k),$$

where $Y_i \stackrel{i.i.d}{\sim} \text{Lap}(\frac{\Delta f}{\epsilon})$.

- Let us apply it in our initial problem where $f(X) = \frac{1}{n} \sum_{i=1}^n X_i$ and $X_i \in \{0, 1\}$.
- We have $\Delta f = 1/n$, therefore $\tilde{p} = f(X) + Y$ where $Y \sim \text{Lap}(\frac{1}{\epsilon n})$.
- Recalling that $p = f(X)$, $\mathbb{E}[\tilde{p}] = p$ and $\text{var}(\tilde{p}) = \text{var}(Y) = \frac{2}{\epsilon^2 n^2}$, we get:

$$|\tilde{p} - p| = O_p\left(\frac{1}{\epsilon n}\right),$$

which is quadratically smaller than the accuracy of the Randomized Response algorithm.

- Notice that while $p \in [0, 1]$, $\tilde{p} \in \mathbb{R}$.

Accuracy and ε -DP of the Laplace Mechanism

The Laplace Mechanism is ε -DP

Let X and X' be two neighbouring datasets, and $p_X(z)$ and $p_{X'}(z)$ the pdfs of $\mathcal{M}(X)$ and $\mathcal{M}(X')$ evaluated at an arbitrary point $z \in \mathbb{R}^k$. Then

$$\frac{p_X(z)}{p_{X'}(z)} \leq e^\varepsilon.$$

Accuracy of the Laplace Mechanism

Let $f: \chi^n \rightarrow \mathbb{R}^k$, and $\mathcal{M}(X) = f(X) + (Y_1, \dots, Y_k)$, where $Y_i \stackrel{i.i.d}{\sim} \text{Lap}(\frac{\Delta f}{\varepsilon})$. Then $\forall \beta \in (0, 1]$

$$\Pr \left(\|f(X)_i - \mathcal{M}(X)_i\|_\infty \geq \ln \left(\frac{k}{\beta} \right) \left(\frac{\Delta f}{\varepsilon} \right) \right) = \beta$$

which yields

$$|f(X) - \mathcal{M}(x)| = O_p \left(\frac{\Delta f}{\varepsilon} \right).$$

Inference in the DP Framework

- Given a suitable (sufficient) statistic that we privatize, we have three ways to make inference:
 - 1 Determine the **exact distribution** of the privatized statistic under the null.
 - 2 Use a **suitable approximation** of its distribution under the null.
 - 3 Use the **non-private distribution** under the null.
- To determine an exact distribution is not always trivial: the convolutions of the added noise and the sampling error is usually intractable.
- To use an approximate distribution or the non-private one is more straightforward but is often far from accurate in finite samples.

A Simple Proportion

Framework

Let $U_{0,j} \stackrel{i.i.d}{\sim} \text{Unif}(0,1)$ be a given seed with $j = 1, \dots, n$.

From the vector of seeds $\mathbf{U}_0 = (U_{0,1}, \dots, U_{0,n})$, a sample is generated as follows:

$$Z(\theta_0, \mathbf{U}_0, n) = (\underbrace{\mathbb{1}_{\{U_{0,1} < \theta_0\}}}_{Z_{0,1}}, \dots, \underbrace{\mathbb{1}_{\{U_{0,n} < \theta_0\}}}_{Z_{0,n}}), \quad Z_{0,j} \stackrel{i.i.d}{\sim} \text{Ber}(\theta_0).$$

Let

$$X_0 := X(\theta_0, \mathbf{U}_0, n) = \sum_{j=1}^n Z_{0,j},$$

from which we compute the sample proportion

$$\hat{\theta}_0 := \hat{\theta}(\theta_0, \mathbf{U}_0, n) = \frac{X(\theta_0, \mathbf{U}_0, n)}{n}.$$

Private Inference approaches

Let us consider the following three approaches:

DP-UMP Test (*Awan and Slavković, 2018*)

Consider $Y_0 = X_0 + N_0$ where $N_0 \sim \text{Tulap}(0, e^{-\varepsilon}, 0)$, then $Y_0|X_0 \sim \text{Tulap}(X_0, e^{-\varepsilon}, 0)$.

- Releasing $Y_0|X_0$ satisfies ε -DP.
- A UMP size- α test can be derived and p -values can be computed easily.

DP Normal Approximation Test (*Vu and Slavković, 2009*)

Consider $\hat{\theta}_0^{(p)} = \hat{\theta}_0 + L_0$ where $L_0 \sim \text{Lap}(\frac{1}{\varepsilon n})$. The distribution of $\hat{\theta}_0$ can be approximated with $N(\theta_0, \frac{\sigma^2}{n})$ and if L_0 is approximated with a $N(0, \frac{2}{\varepsilon^2 n^2})$, then under the null

$$\hat{\theta}_0^{(p)} \sim N\left(\theta_0, \frac{\sigma^2}{n} + \frac{2}{\varepsilon^2 n^2}\right).$$

Non-private JIMI

Non-private JIMI Estimator

The **Just Identified Minimal distance estimator (JIMI)** in the non-private framework is defined as follows:

Let $U_{ij} \stackrel{i.i.d}{\sim} \text{Unif}(0, 1)$ be a given seed with $i = 1, \dots, B$ and $j = 1, \dots, n$.

$$\tilde{\theta}_i \in \underset{\theta}{\text{argzero}} \hat{\theta}(\theta_0, \mathbf{U}_0, \varepsilon, n) - \hat{\theta}(\theta, \mathbf{U}_i, \varepsilon, n),$$

which implies in the case of a proportion

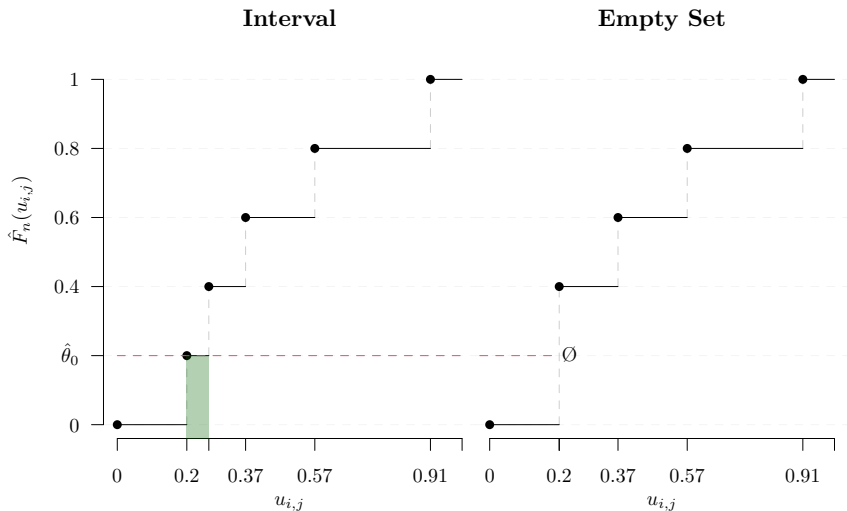
$$\tilde{\theta}_i \in \left\{ \theta : \frac{\sum_{j=1}^n \mathbb{1}\{U_{ij} < \theta\}}{n} = \hat{\theta}(\theta_0, \mathbf{U}_0, n) \right\} \Leftrightarrow \tilde{\theta}_i \in \left\{ \theta : \hat{F}_n(\theta; \mathbf{U}_i) = \hat{\theta}(\theta_0, \mathbf{U}_0, n) \right\}.$$

This in turn implies that, when non-empty, the solution is an interval

$$\tilde{\theta}_i \in \hat{F}_n^-(\hat{\theta}_0; \mathbf{U}_i) = [u_{i, [n\hat{\theta}_0]}, u_{i, [n\hat{\theta}_0+1]}),$$

for $\hat{\theta}_0 \in (0, 1)$, otherwise take $[0, u_{i, [1]}]$ if $\hat{\theta}_0 = 0$ and $[u_{i, [n]}, 1]$ if $\hat{\theta}_0 = 1$, where $q_{[x]}$ corresponds to the order statistic of rank x .

Solution Types



Selection Criterion

Selection Criterion

In order to have a unique solution, we need to impose a selection rule on $\tilde{\theta}_i$. Let D be a selection criterion:

$$\tilde{\theta}_i = Q(\hat{\theta}_0, D, \mathbf{U}_i) = U_{i, [n\hat{\theta}_0]} + D \left(U_{i, [n\hat{\theta}_0+1]} - U_{i, [n\hat{\theta}_0]} \right). \quad (1)$$

It turns out that the exact distribution of $\tilde{\theta}_i$ can sometimes be derived depending on the choice of D as demonstrated by [Hannig \(2009\)](#). Particularly,

Distribution of D	Distribution of $\tilde{\theta}_i$
Beta $(1/2, 1/2)$	Beta $(n\hat{\theta}_0 + 1/2, n(1 - \hat{\theta}_0) + 1/2)$
Beta $(1, 1)$	Beta $(n\hat{\theta}_0 + 1, n(1 - \hat{\theta}_0) + 1)$

Therefore, in stead of solving numerically, simply use the exact distribution when available.

Differentially Private JIMI

DP Proportion

Consider a simple DP proportion

$$\hat{\pi}(\theta_0, \mathbf{U}_0, W_0, \varepsilon, n) := \frac{X(\theta_0, \mathbf{U}_0, n)}{n} + Y(W_0, \varepsilon, n),$$

where

$$Y(W_0, \varepsilon, n) := -\frac{1}{\varepsilon n} \text{sign}(W_0) \log(1 - |W_0|)$$

with $W_i \sim \text{Unif}(-1/2, 1/2)$. Thus

$$Y(W_0, \varepsilon, n) \sim \text{Laplace}\left(0, (\varepsilon n)^{-1}\right)$$

Differentially Private JIMI

DP JIMI

We define the DP JIMI estimator as

$$\tilde{\theta}_i \in \underset{\theta}{\operatorname{argzero}} \hat{\pi}(\theta_0, \mathbf{U}_0, W_0, \varepsilon, n) - \hat{\pi}(\theta, \mathbf{U}_i, W_i, \varepsilon, n),$$

with $i \in 1, \dots, B$, which implies

$$\tilde{\theta}_i \in \{\theta : X(\theta, \mathbf{U}_i, n) = [\hat{\pi}(\theta_0, \mathbf{U}_0, W_0, \varepsilon, n) - Y(W_i, \varepsilon, n)] n\}$$

However, since $X(\theta, \mathbf{U}_i, n) \in \{0, 1, \dots, n\}$, the solution interval is rarely non-empty. Therefore, let us instead solve

$$\tilde{\theta}_i^* \in \underset{\theta}{\operatorname{argmin}} \hat{\pi}(\theta_0, \mathbf{U}_0, W_0, \varepsilon, n) - \hat{\pi}(\theta, \mathbf{U}_i, W_i, \varepsilon, n),$$

and we can write

$$\tilde{\theta}_i^* = Q^*(\hat{\theta}_i^*, D, \mathbf{U}_i) = U_{i, \lfloor n\hat{\theta}_i^* \rfloor} + D \left(U_{i, \lfloor n\hat{\theta}_i^* \rfloor + 1} - U_{i, \lfloor n\hat{\theta}_i^* \rfloor} \right), \quad (2)$$

where $\hat{\theta}_i^* = \hat{\pi}(\theta_0, \mathbf{U}_0, W_0, \varepsilon, n) - Y(W_i, \varepsilon, n)$, and $q_{\lfloor x \rfloor}$ returns the order statistic whose rank is the closest integer to x .

DP JIMI algorithm

Using the same rule as before, we get

$$\tilde{\theta}_i^* = Q^*(\hat{\theta}_i^*, D, \mathbf{U}_i) \xrightarrow{D \sim \text{Beta}(1/2, 1/2)} \tilde{\theta}_i^* \sim \text{Beta}(n\hat{\theta}_i^* + 1/2, n(1 - \hat{\theta}_i^*) + 1/2). \quad (3)$$

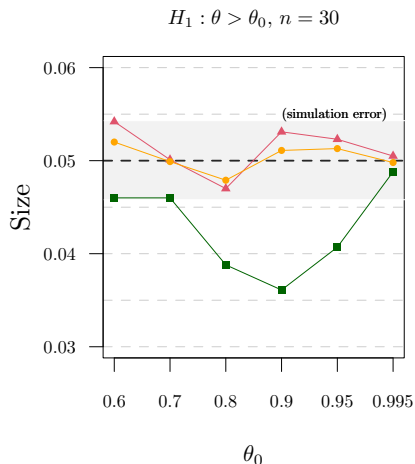
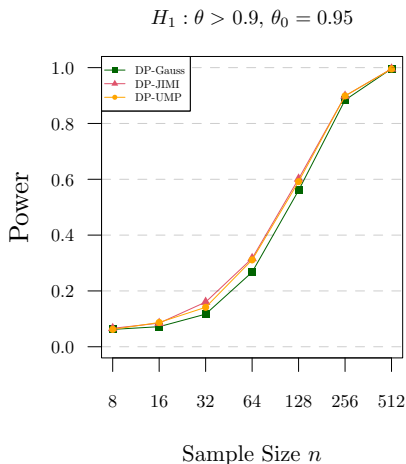
Algorithm:

INPUT: $n \in \mathbb{N}$, $\varepsilon > 0$, $W_0 \sim \text{Unif}(-1/2, 1/2)$, $\mathbf{U}_0 \sim \text{Unif}(0, 1)$, $\theta_0 \in (0, 1)$,
 $\hat{\pi}(\theta_0, \mathbf{U}_0, W_0, \varepsilon, n) \in \mathbb{R}$, $i = 1, \dots, B$

- 1 Generate $\hat{\theta}_i^* := \hat{\pi}_0 - Y_i$
- 2 if $\hat{\theta}_i^* > 1$ return $\tilde{\theta}_i^* = 1$,
if $\hat{\theta}_i^* < 0$ return $\tilde{\theta}_i^* = 0$,
otherwise, generate $\tilde{\theta}_i^*$ from $\text{Beta}(n\hat{\theta}_i^* + 1/2, n(1 - \hat{\theta}_i^*) + 1/2)$

OUTPUT: A differentially private JIMI estimate $\tilde{\theta}_i^* \in [0, 1]$

One-sample Proportion Tests: Power and Size



First Order Properties

- Given $n \geq 1$, $M_{i,j}, R_{i,j} \sim \Gamma\left(\frac{1}{n}, \frac{1}{\varepsilon n}\right)$, then $\sum_{j=1}^n (M_{i,j} - R_{i,j}) \sim \text{Lap}\left(\frac{1}{\varepsilon n}\right)$.
- Hence, the DP proportion can be written as

$$\hat{\pi}_0 = \frac{1}{n} \sum_{j=1}^n (Z_{0,j} + n(M_j - R_j)).$$

- Let $V_{0,j} := Z_{0,j} + n(M_{0,j} - R_{0,j})$, then the MGF of $V_{0,j}$ is

$$M_{V_{0,j}}(t) = \frac{1 - \theta_0 + \theta_0 e^t}{1 - \varepsilon^{-2} t^2}, \quad |t| < \varepsilon,$$

and as $M_{V_{0,j}}(t)$ is well defined in $[-\varepsilon, \varepsilon]$ with $\varepsilon > 0$, all the moments are finite.

- Then by the Berry-Esseen theorem, it follows that

$$\sqrt{n} \sigma(\theta)^{-1} (\hat{\pi}_n(\theta) - \theta) \rightarrow \mathbf{Z}, \text{ uniformly in } \theta \in (0, 1) \text{ (Assumption 2),}$$

where $\sigma^2(\theta) = \theta(1 - \theta) + 2/\varepsilon^2$. This in turn implies that

$$\sqrt{n} \sigma(\hat{\pi}_n)^{-1} (\tilde{\theta}_n^* - \hat{\pi}_n) \rightarrow \mathbf{Z}, \text{ in probability (Assumption 1),}$$

and ensures that the DP-JIMI confidence intervals are asymptotically consistent (see Lemma 23.3, in [Van der Vaart \(2000\)](#))

Conclusions, Limitations and Possible Extensions...

- Recall that $|f(X) - \mathcal{M}(x)| = O_p\left(\frac{\Delta f}{\varepsilon}\right)$, what if Δf is not bounded?
 - Before adding the noise, the curator will truncate/clip the data to $[a, b] \subset \mathbb{R}$ before computing $f(X)$.
 - In this case the non-private initial estimator $f(X)$ will **no longer be consistent**, and inference becomes problematic for many methods.
 - This is much less of an issue for the JIMI.
- Our method remains simple to apply (no calculations required), provided we know how to generate the data. What if the covariates are privatized?
 - We find ourselves in an **error of measurement** setting, which remains at the moment a limitation.

Selected References

- Jordan Awan and Aleksandra Slavković. Differentially private uniformly most powerful tests for binomial data. *Advances in Neural Information Processing Systems*, 31, 2018.
- Cynthia Dwork, Aaron Roth, et al. The algorithmic foundations of differential privacy. *Foundations and Trends® in Theoretical Computer Science*, 9(3–4): 211–407, 2014.
- Jan Hannig. On generalized fiducial inference. *Statistica Sinica*, pages 491–544, 2009.
- Aad W Van der Vaart. *Asymptotic statistics*, volume 3. Cambridge university press, 2000.
- Duy Vu and Aleksandra Slavkovic. Differential privacy for clinical trial data: Preliminary evaluations. In *2009 IEEE International Conference on Data Mining Workshops*, pages 138–143. IEEE, 2009.